

**DR. BABASAHEB AMBEDKAR TECHNOLOGICAL UNIVERSITY, LONERE –
RAIGAD -402 103
Mid Semester Examination – October - 2017**

Branch: M.Tech VLSI

Sem.: - I

Subject with Subject Code: - Graph Theory and Discrete Optimization (MTVLC101)

Marks: 20

Date: -09-10-2017

Time: - 1 Hr.

MODEL ANSWER SOLUTION

Instructions: - Assume necessary data wherever required.

(Marks)

Q.No.1 Attempt any one of the following

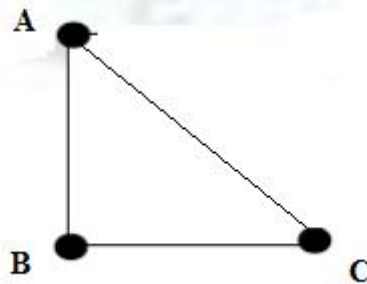
(08)

a.) What are Hamiltonian circuit, Hamiltonian trail and Hamiltonian graph? Prove that if G fails Chavatal's condition, then G has at least $(n - 2)$ edges. Conclude from this that the maximum number of edges in a simple non-Hamiltonian n -vertex graph is $\binom{n-1}{2} + 1$.

Ans: Hamiltonian trail of a graph is a one which visits every vertex only once and it can start from any vertex and end to any vertex.

Hamiltonian circuit is a closed path that visits every vertex only once, start and end vertex is same.

A connected graph G is said to be a Hamiltonian graph, if there exists a cycle which contains all the vertices of G .



Example:

For the above graph the Hamiltonian trail is $A \rightarrow B \rightarrow C \rightarrow A$.

As the trail contains all the vertices it is Hamiltonian graph. The graph start and end on the same vertex so it is Hamiltonian circuit.

Suppose G fails Chavatal's condition. Then there exists some $i < n/2$ such that $d_i \geq i$ and $d_{n-i} < n - i$. Let u be a vertex with degree d_i , and let v be a vertex with degree d_{n-i} . Thus, in G , we have

$$d_{\bar{G}}(u) + d_{\bar{G}}(v) = (n - 1) - d_G(u) + (n - 1) - d_G(v)$$

$$\begin{aligned}
 &= 2n - 2 - d_i - d_{n-i} \\
 &> 2n - 2 - i - (n - i) = n - 2
 \end{aligned}$$

So

$$d_{\bar{G}}(u) + d_{\bar{G}}(v) \geq n - 1$$

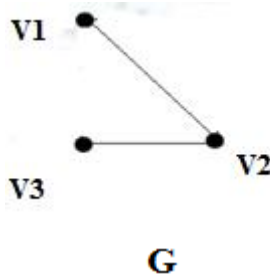
Since u and v have degree sum at least $n - 1$ in \bar{G} , and since a simple graph has at most one edge joining them (counted twice in the degree sum), there must be at least $n - 2$ distinct edges in \bar{G} incident to $\{u, v\}$. So $|E(\bar{G})| \geq n - 2$.

Therefore, the number of edges in a simple non-Hamiltonian graph G on n vertices is

$$\begin{aligned}
 |E(G)| &= \binom{n}{2} - |E(\bar{G})| \leq \frac{n(n-1)}{2} - (n-2) \\
 &= \frac{n^2 - 3n + 4}{2} = \frac{(n-1)(n-2)}{2} + 1 = \binom{n-1}{2} + 1
 \end{aligned}$$

So the maximum number of edges is $\binom{n-1}{2} + 1$.

b.) What is a spanning tree give example? Find the number of spanning trees for the below incomplete graph G .

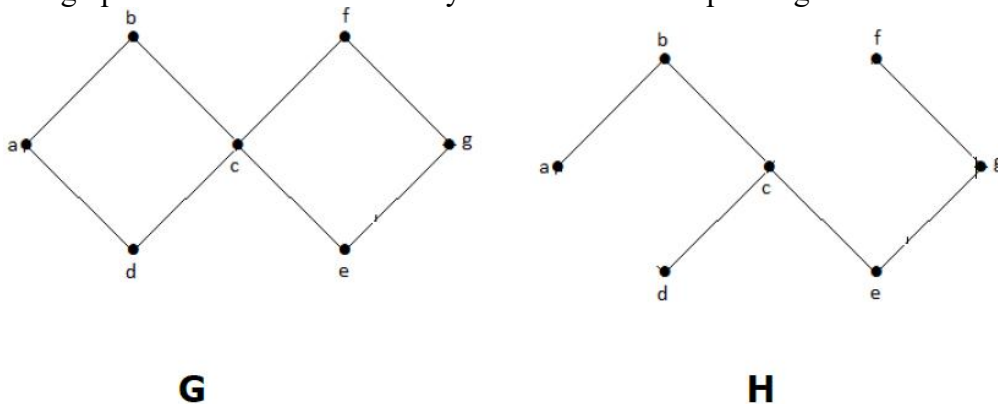


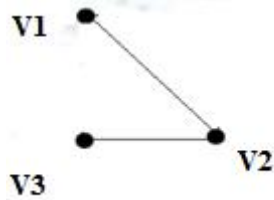
Ans:

A spanning tree of an undirected graph is the sub-graph that includes all the vertices of G .

Example: In the below example, G is the connected graph and H is the sub-graph of G .

The graph H is a tree as it has no cycles. Hence it is a spanning tree.





Here to find the spanning tree.

1. Construct the adjacency matrix.

$$\begin{bmatrix} 0 & 0 & 1 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \end{bmatrix}$$

2. Diagonal 0's should be replaced with degree of node.

$$\begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 1 & 1 & 2 \end{bmatrix}$$

3. Non-diagonal 1's should be replaced with "-1".
4. Non-diagonal 0's will remain 0's.

$$\begin{bmatrix} 1 & 0 & -1 \\ 0 & 1 & -1 \\ -1 & -1 & 2 \end{bmatrix}$$

5. No of spanning trees = co-factor of any element.

$$a_{11} = (-1)^{1+1} \begin{vmatrix} 1 & -1 \\ -1 & 2 \end{vmatrix}$$

Therefore, No of spanning tree = 1

Q.No. 2 Attempt any three of the following:

(12)

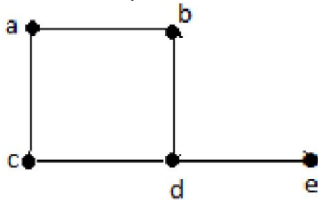
a.) Define. (Give an example)

1. Graph

Ans: A graph is a pair of sets (V, E), where V is the set of vertices and E is the set of edges, connecting the pair of vertices. In the following graph

$$V = \{a, b, c, d, e\}$$

$$E = \{ab, ac, bd, bc, de\}$$

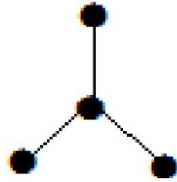


2. Tree

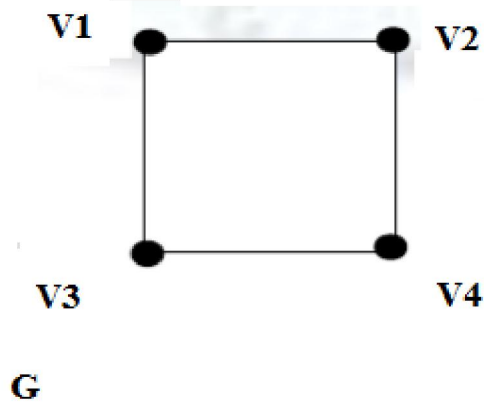
Ans: A connected acyclic graph is called a tree. In other words, a connected graph with no cycles is called a tree.

The edges of a tree are known as **branches**. Elements of trees are called their **nodes**. The nodes without child nodes are called **leaf nodes**.

Example: The graph shown here is a tree because it has no cycles and it is connected. It has four vertices and three edges, i.e., for 'n' vertices 'n-1' edges.



b.) What is an Adjacency matrix and Incidence matrix of graph? Find the adjacency matrix of G shown in figure below



Ans:

An adjacency matrix is a way to represent a finite graph G on n vertices by an nxn matrix a, whose entries are either 0 or 1.

Consider a graph G having vertex set $V(G) = \{V_1, V_2, \dots, V_n\}$. Then the adjacency matrix $A(G) = [a_{ij}]$ of G is a n x n matrix defined as

$$a_{ij} = \begin{cases} 1 & \text{if } v_i \text{ is adjacent to } v_j \\ 0 & \text{otherwise} \end{cases}$$

The incidence matrix of a graph is a matrix representation of the graph, in which each row and each column represents an edge and a vertex, respectively.

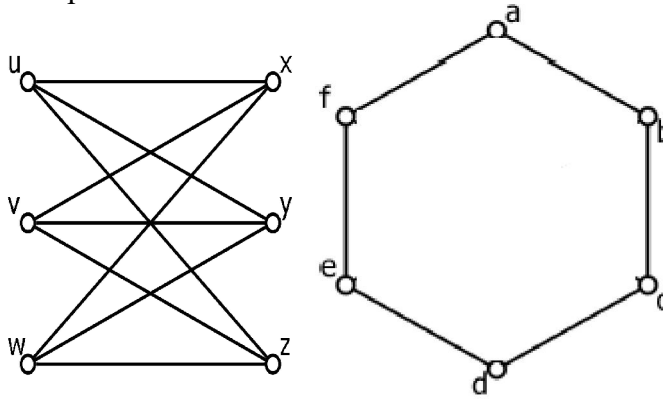
Consider a graph G having vertex set $V(G) = \{v_1, v_2, \dots, v_n\}$ and $E(G) = \{e_1, e_2, \dots, e_m\}$. Then the incident matrix $B = [b_{ij}]$ of G is a n x m matrix defined as

$$b_{ij} = \begin{cases} 1 & \text{if } v_i \text{ and } e_j \text{ are incident} \\ 0 & \text{otherwise} \end{cases}$$

The adjacency matrix of G is

$$A(G) = \begin{matrix} & \begin{matrix} V_1 & V_2 & V_3 & V_4 \end{matrix} \\ \begin{matrix} V_1 \\ V_2 \\ V_3 \\ V_4 \end{matrix} & \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix} \end{matrix}$$

c.) What is an Isomorphic graph? Prove whether the graphs G1 and G2 are isomorphic or not?



G1 G2

Ans:

Two graphs G1 and G2 are said to be isomorphic if:

1. Their numbers of components (vertices and edges) are same.
2. Their edge connectivity is retained.

If $G1 \cong G2$ then:

1. $|V(G1)| = |V(G2)|$
2. $|E(G1)| = |E(G2)|$
3. Degree sequences of G1 and G2 are same.

For the above two graphs G1 and G2

1. Number of vertices are same i.e 6
2. Number of edges for G1 are 9 and for G2 are 6 which are not same.

Therefore, graphs G1 and G2 are non-isomorphic graphs as all the condition for isomorphic graphs are not satisfied.

d.) Define Eccentricity, radius and diameter of graph. Prove that For any connected graph G, $\text{rad}(G) \leq \text{diam}(G) \leq 2 \text{ rad}(G)$.

Ans:

For a given vertex v of a connected graph, the *eccentricity* of v, denoted $\text{ecc}(v)$, is defined to be the greatest distance from v to any other vertex. That is,

$$\text{ecc}(v) = \max\{d(u, v)\}$$

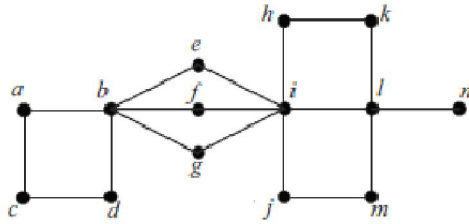
$$x \in V(G)$$

In Figure below, $\text{ecc}(a) = 5$ since the farthest vertices from a (namely k, m, n) are at a distance of 5 from a. Of the vertices in this graph, vertices c, k, m and n have the greatest eccentricity (6), and vertices e, f and g have the smallest eccentricity (3). These values and types of vertices are given special names.

In a connected graph G, the radius of G, denoted $\text{rad}(G)$, is the value of the smallest eccentricity.

The diameter of G, denoted $\text{diam}(G)$, is the value of the greatest eccentricity.

In Figure below, the radius is 3, the diameter is 6.



Proof:

By definition, $\text{rad}(G) \leq \text{diam}(G)$, so we just need to prove the second inequality. Let u and v be vertices in G such that $d(u, v) = \text{diam}(G)$. Further, let c be a vertex in the center of G . Then,

$$\text{diam}(G) = d(u, v) \leq d(u, c) + d(c, v) \leq 2 \text{ecc}(c) = 2 \text{rad}(G)$$